

## FINDING A VALUE ON PARTIALLY DEFINED GAMES

## Introduction

A cooperative game is a pair  $(N, v)$  where  $N = \{1, 2, \dots, n\}$  and  $v$  is a real-valued function on the nonempty subsets of  $N$ . Elements of  $N$  are often called players while subsets of  $N$  are called coalitions. The grand coalition, denoted as  $N$ , is composed when all of the players cooperate together. The function  $v$  is called the worth function, and  $v(S)$  is interpreted as the worth of the coalition  $S$ . In other words,  $v(S)$  is the amount that individuals in  $S$  can jointly obtain if they cooperate as a group.

A game  $(N, v)$  is said to be superadditive if  $v(S \cup T) \geq v(S) + v(T)$  for all disjoint coalitions  $S, T \subseteq N$ . The game  $(N, v)$  is monotonic if  $v(S) \leq v(T)$  for all coalitions  $S \subseteq T \subseteq N$ . The game  $(N, v)$  is said to be 0-normalized if  $v(i) = 0$  for all  $i \in N$ . The 0-normalization of a game  $(N, v)$  is the game  $(N, u)$  where

$$u(S) = v(S) - \sum_{i \in S} v(i).$$

A game  $(N, v)$  is 0-monotonic if its 0-normalization is monotonic.

In an effort to find some standard of fairness or a predictor of bargaining solutions, many researchers of game theory have concentrated on finding the "best" way for individuals in a game to form coalitions and eventually maximize their savings. Once the savings have been made they must be allocated to the participating players. An allocation method is a function that assigns an

allocation to each cooperative game in some class.

For every cooperative game there are  $2^N - 1$  possible coalitions. Normally the worth of every coalition is known. However, this may not always be the case. For instance, a firm might have certain time or monetary constraints that prevent it from knowing the worth of all coalitions. The utility for a game of this type occurs when it is impractical to determine every coalitional worth,  $v(S)$ . During the summer of 1990, while involved in a Research Experience for Undergraduates program at Drew University, David Letscher developed and began to research what might happen in such cases. He named this topic partially defined games.

#### Necessary definitions

A partially defined game, or PDG, is a triple  $(N, Z, v)$  where  $N = \{1, 2, \dots, n\}$  is the set of players,  $Z$  is a collection of nonempty subsets of  $N$ , and  $v$  is a real-valued function on  $Z$ . We will often use  $v$  to denote a partially defined game when  $N$  and  $Z$  are clear from context.

Definition - Let  $J$  be a subset of the player set  $N$  such that  $1, n \in J$ . A J-game is a partially defined game where  $Z = \{S \subseteq N \mid |S| \in J\}$ .

We will denote a J-game by  $(N, J, v)$ .  $Z$  is the set of coalitions whose worths are known, while  $J$  is the set of sizes, or cardinalities, of coalitions whose worths are known. For example, in a 4 - player partially defined game where the worths of the grand coalition, the triples, and the singletons are known while the worths of the pairs are not,  $J = \{1, 3, 4\}$ . For this same game,  $Z = \{A, B, C, D, ABC, ABD, ACD, BCD, ABCD\}$ . In our research,

